

Very High-Order Masking: Efficient Implementation and Security Evaluation



European Research Council
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CHES 2017, Taipei, Taiwan

Outline

- **Background**
 - Masking
 - Barthe *et al.* masking scheme
- **How fast can be very high-order masking ?**
 - Data representation
 - AES results and discussion
- **How can we evaluate security at very high order ?**
 - Limitation of leakage detection strategy
 - Multi-model approach
- **Conclusion/Open problems**

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- Masking (e.g. Boolean encoding)

$$a = a_1 \oplus a_2 \oplus \cdots \oplus a_d$$

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- Abstract security
 - Probing model
 - Security order
 $d - 1$ (at best)

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- Concrete security
 - Noisy leakage model
 - $N = (\sigma^2)^{d-1}$ (under assumptions)

- Parallel masking scheme by design
- All shares manipulated at once

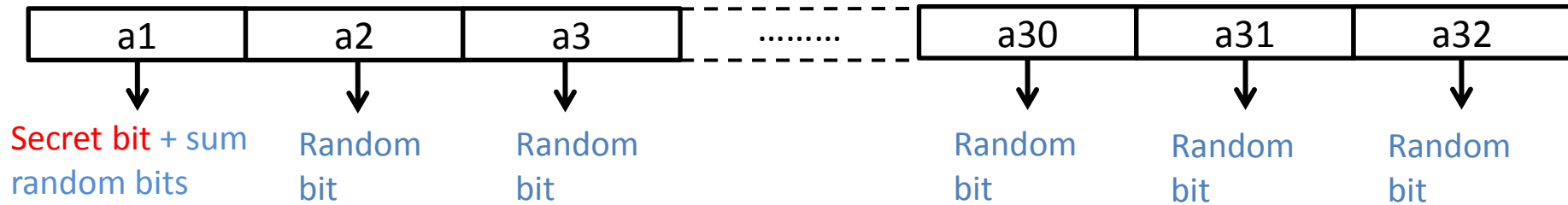
- Parallel masking scheme by design
- All shares manipulated at once
- Example of mult. $a * b = c$ for $d = 3$

$$\begin{array}{c|c|c|c|c|c|c|c|c|c}
 a_1 b_1 & & r_1 & & a_1 b_3 & & a_3 b_1 & & r_3 & & c_1 \\
 a_2 b_2 & \oplus & r_2 & \oplus & a_2 b_1 & \oplus & a_1 b_2 & \oplus & r_1 & \equiv & c_2 \\
 a_3 b_3 & & r_3 & & a_3 b_2 & & a_2 b_3 & & r_2 & & c_3
 \end{array}$$

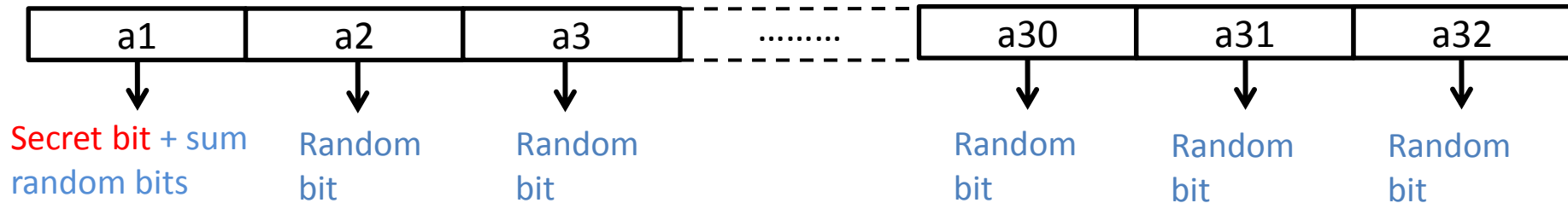
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- 32-bit register

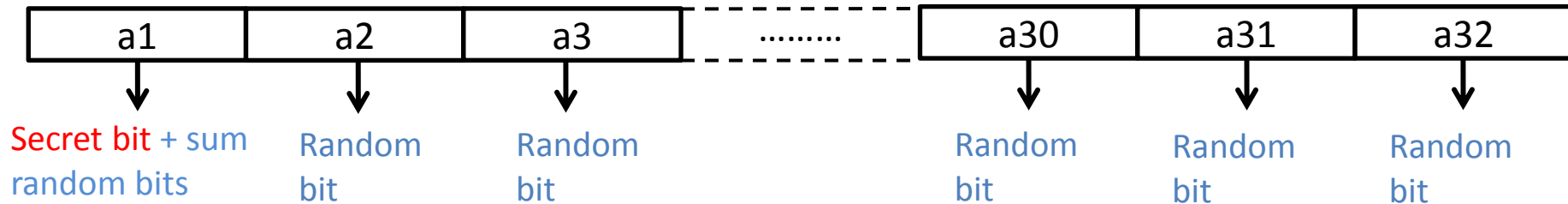


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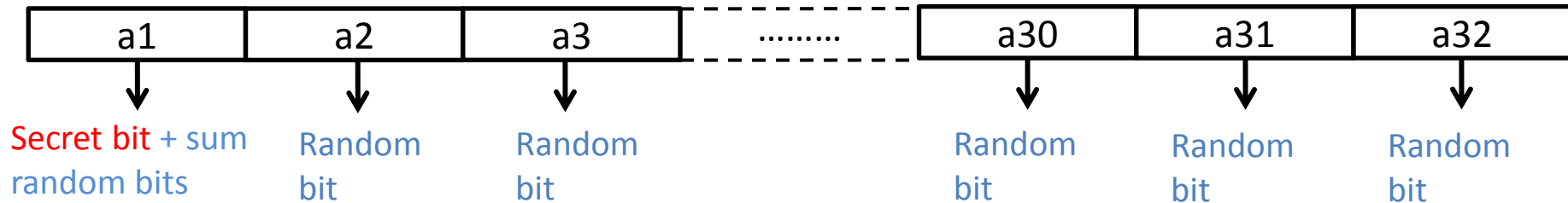
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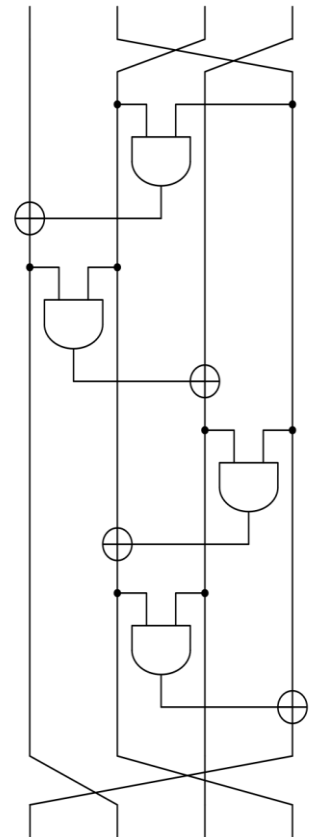


- Use bitwise operators (XOR, AND, ...)
- Implementation on 32-bit ARM
- Optimal case: register size = nb of shares

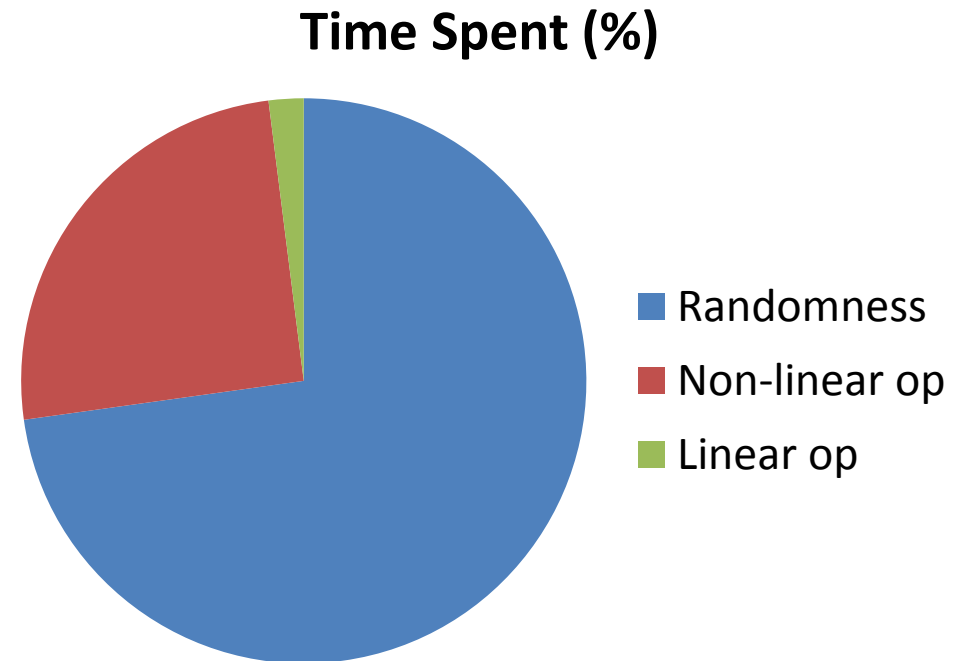
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- Use bitwise operators (XOR, AND, ...)
- Implementation on 32-bit ARM
- Optimal case: register size = nb of shares
- Well suited for bitslice ciphers →



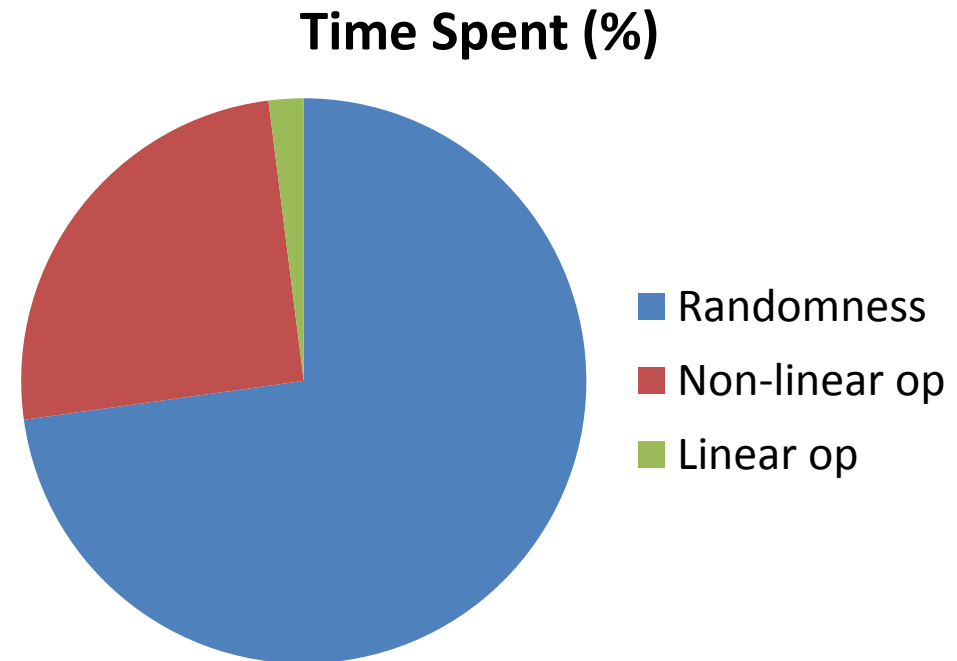
- Application to AES
- Gate level representation of AES S-box (Boyar, Peralta 2010)



10 cycles to generate 32-bit random value

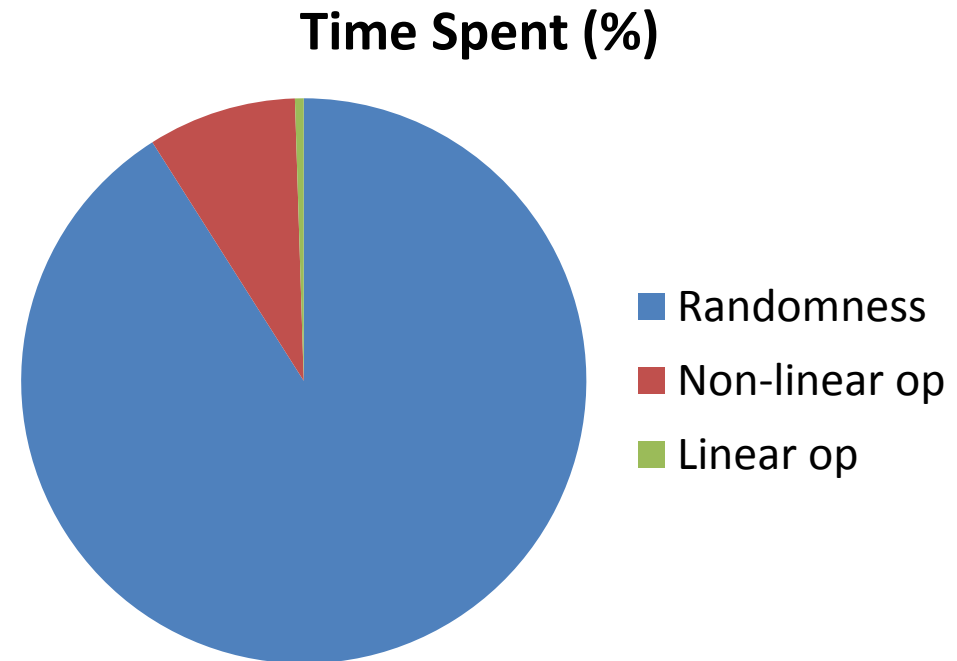
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80 cycles to generate 32-bit random value
Total = 9 700 000 cycles

- Goudarzi-Rivain 2017: Generic ISW implementation and application to bitsliced AES

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3,821,312	2,783,510

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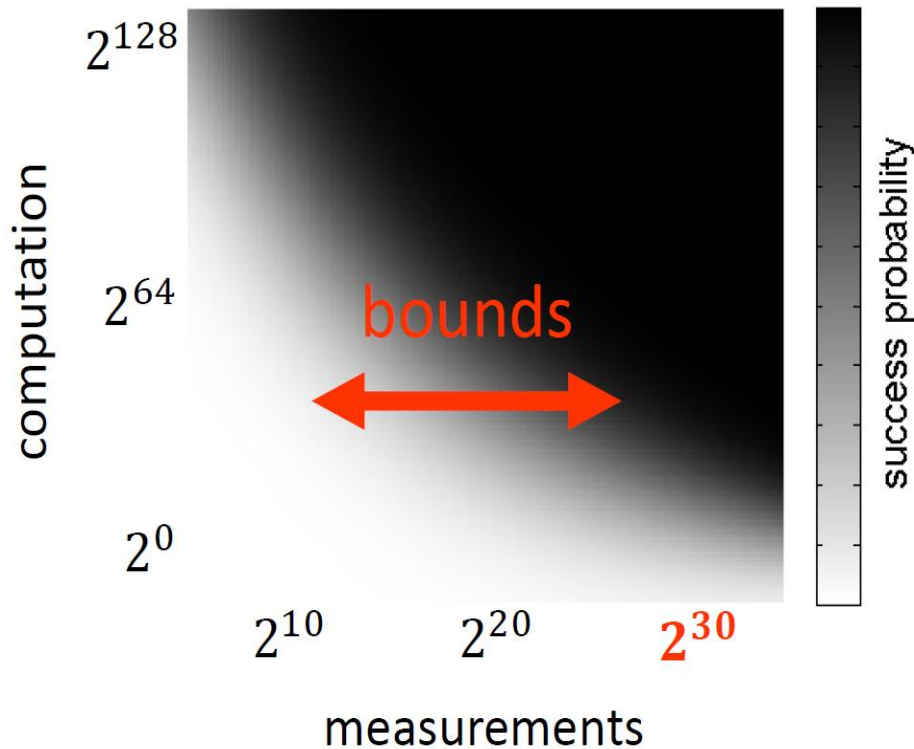
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- Same order of magnitude of cycles
- Very high-order masking is **not out of reach !**

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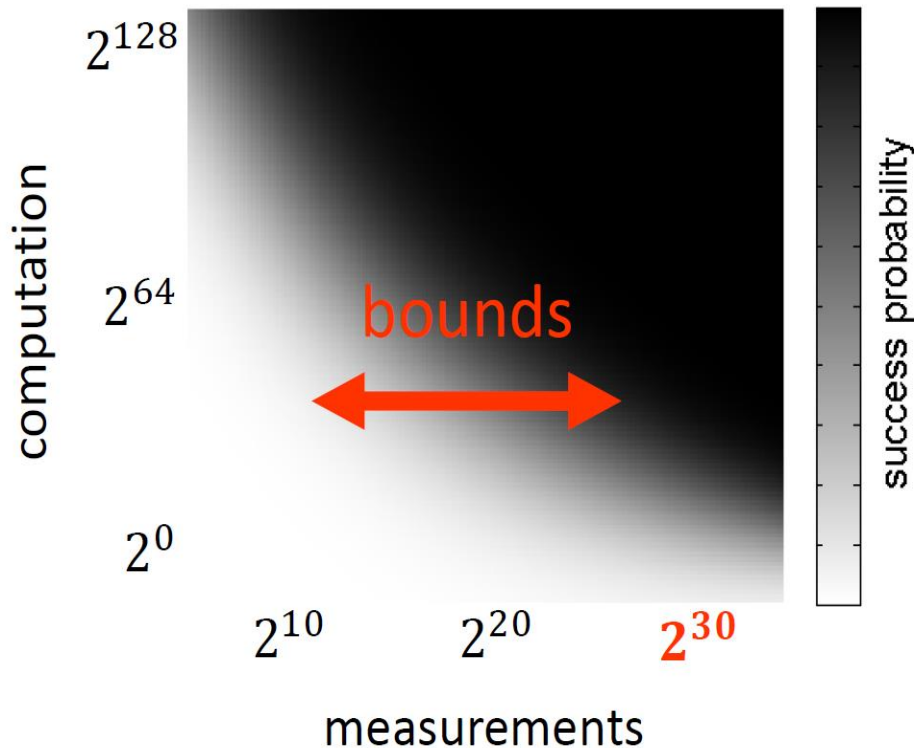
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attack-based evaluations



- Evaluator power = 2^{30}
- If security $\leq 2^{30}$, security level
- What if security $> 2^{30}$?
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We expect 31th-security order (or 31/f-security order)

Probing model

Abstract

Qualitative

Algorithmic security
order
d

Risk captured:

Lack of refreshing

Probing model	Bounded-Moment Model
Abstract	Physical
Qualitative	Qualitative
Algorithmic security order d	Physical security order f
Risk captured: Lack of refreshing	Risk captured: Share recombination

Probing model	Bounded-Moment Model	Noisy Leakage Model
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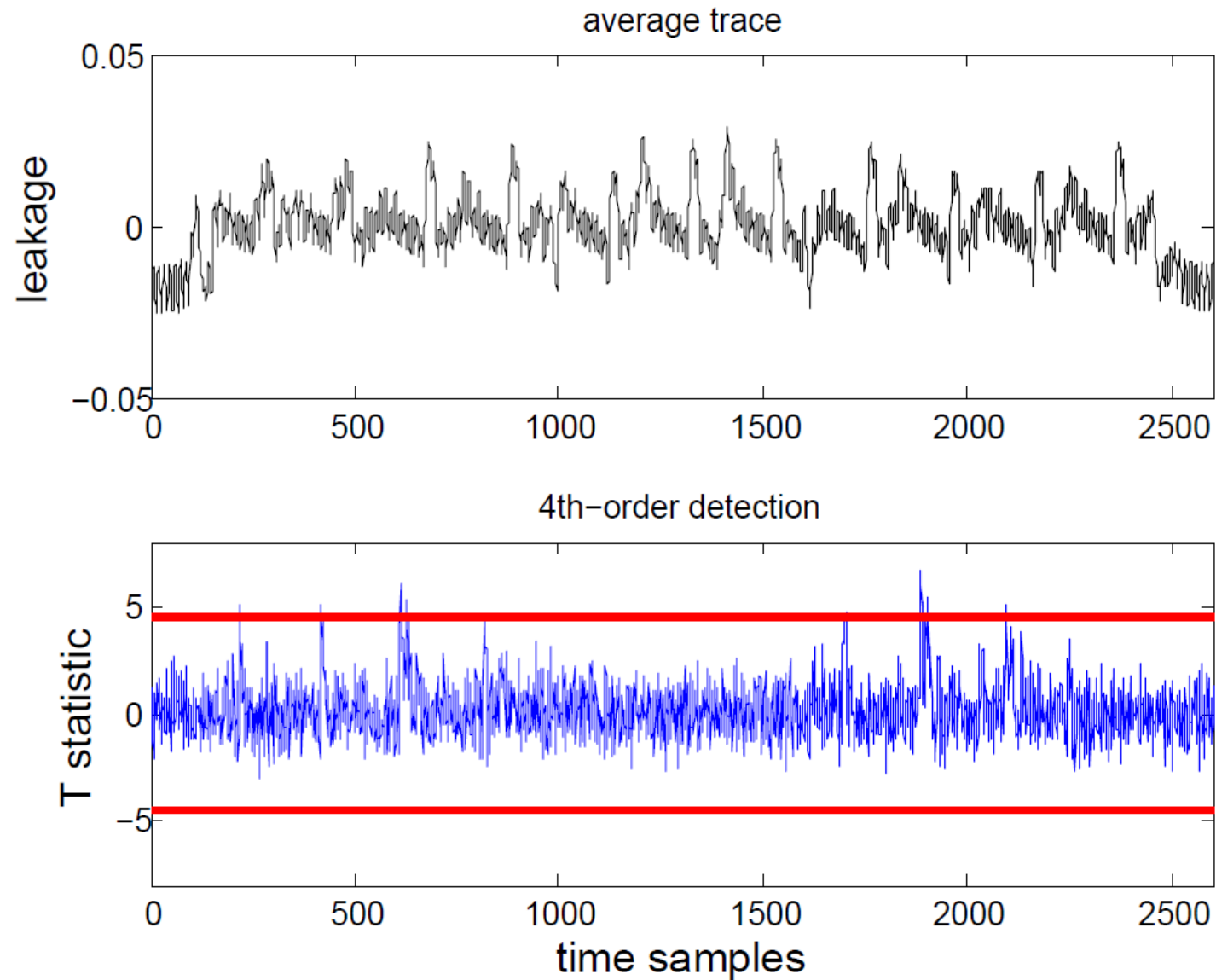
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Abstract	Physical	Physical
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Algorithmic security order	Physical security order	Physical security order
$d + f + \text{SNR} + \text{MI} \Rightarrow \text{Security level}$		
Risk captured: Lack of refreshing	Risk captured: Share recombination	Risk captured: Lack of noise

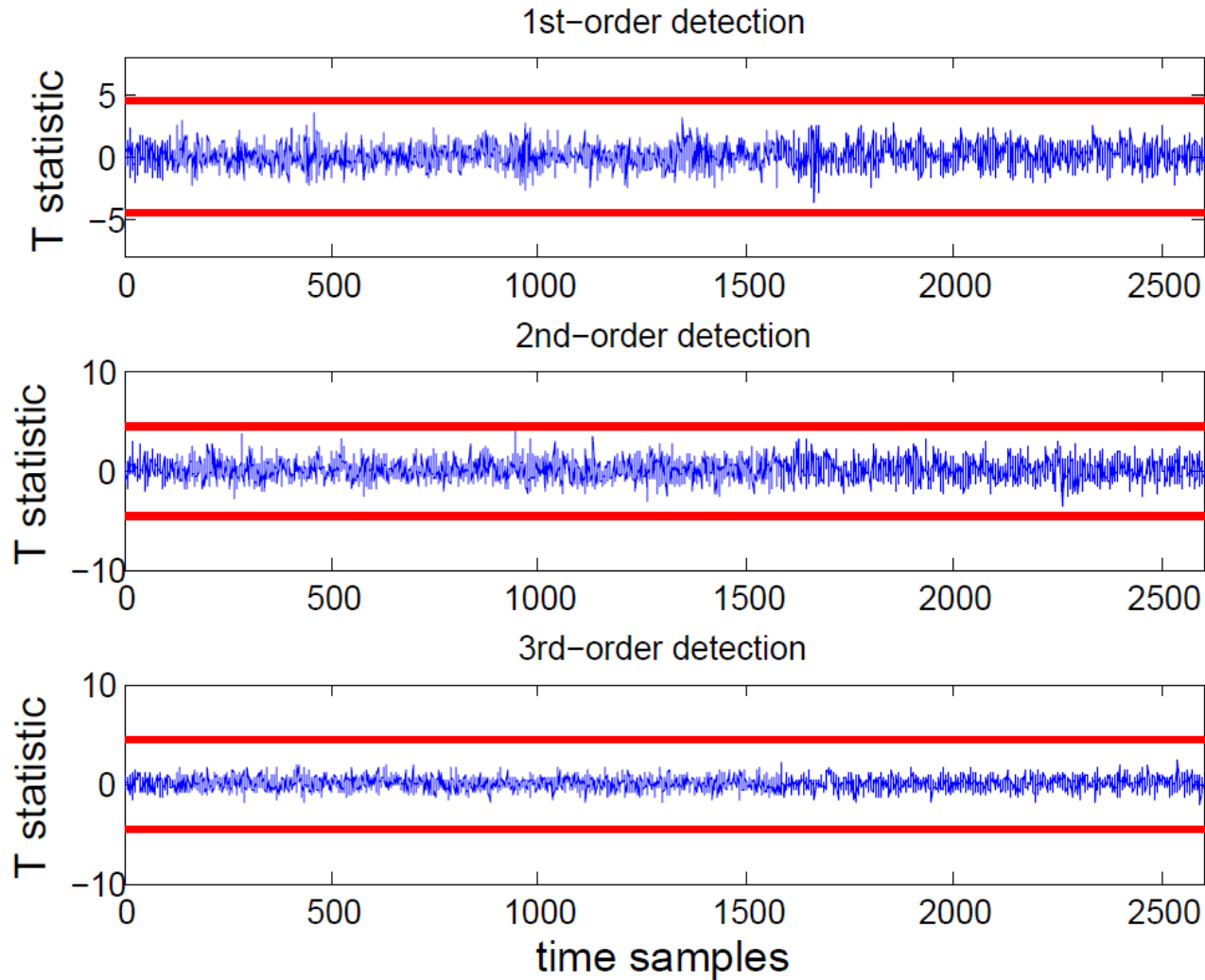
- 2 possible options:
 - Composable gadgets (SNI)
 - Simple to analyse
 - Implementation becomes expensive
 - Full code evaluation
 - Hard to analyse
 - Reduced implementation cost

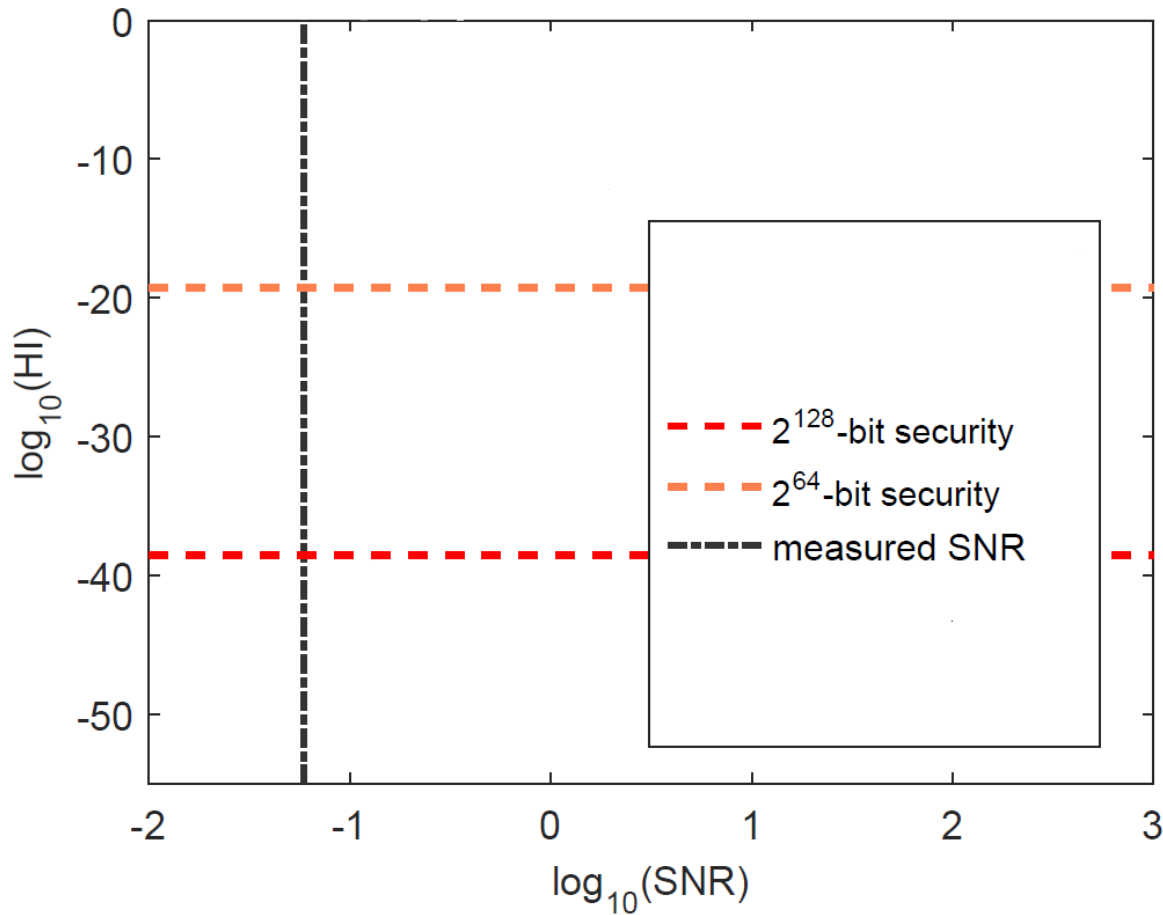
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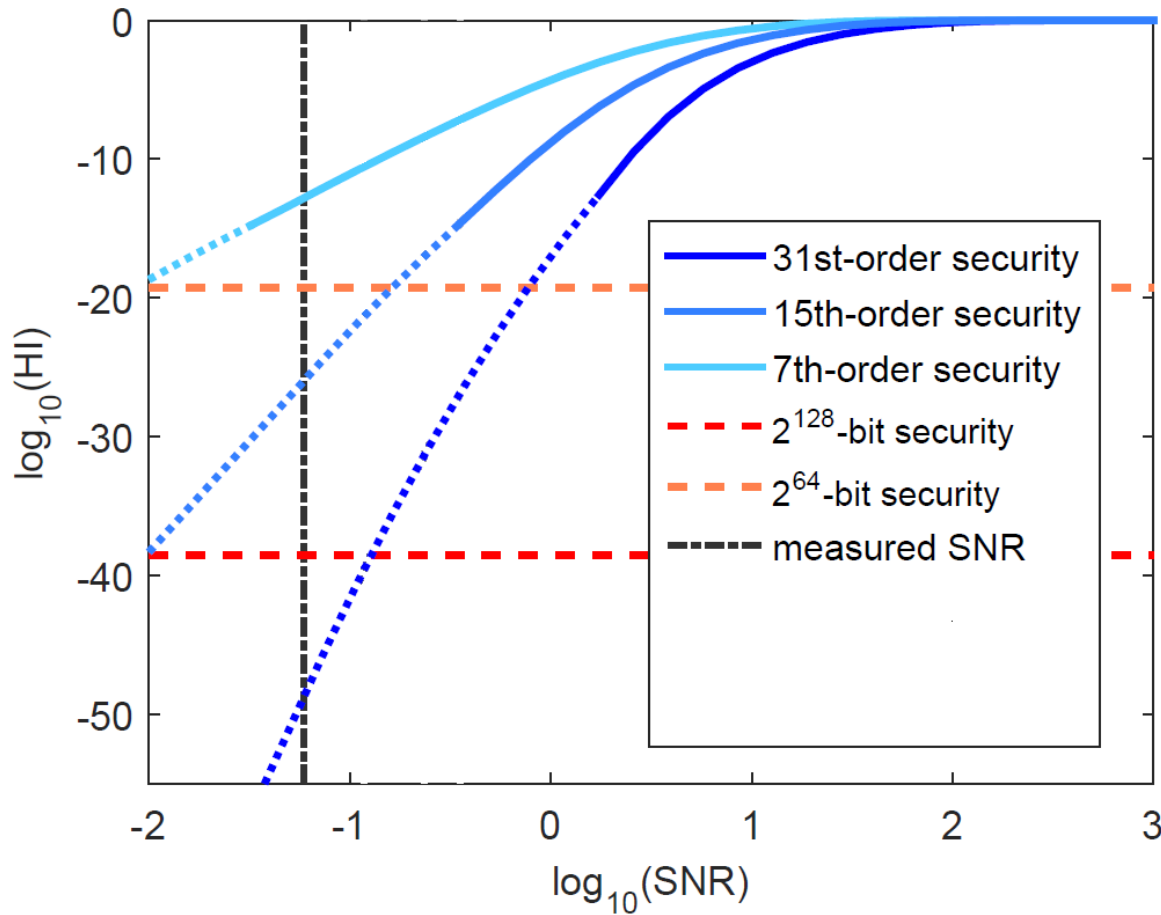
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- Leakage detection on small order (e.g. on 4 shares)
- Extraction of a risk factor \mathbf{f} from possible share recombination
- Extrapolation of security



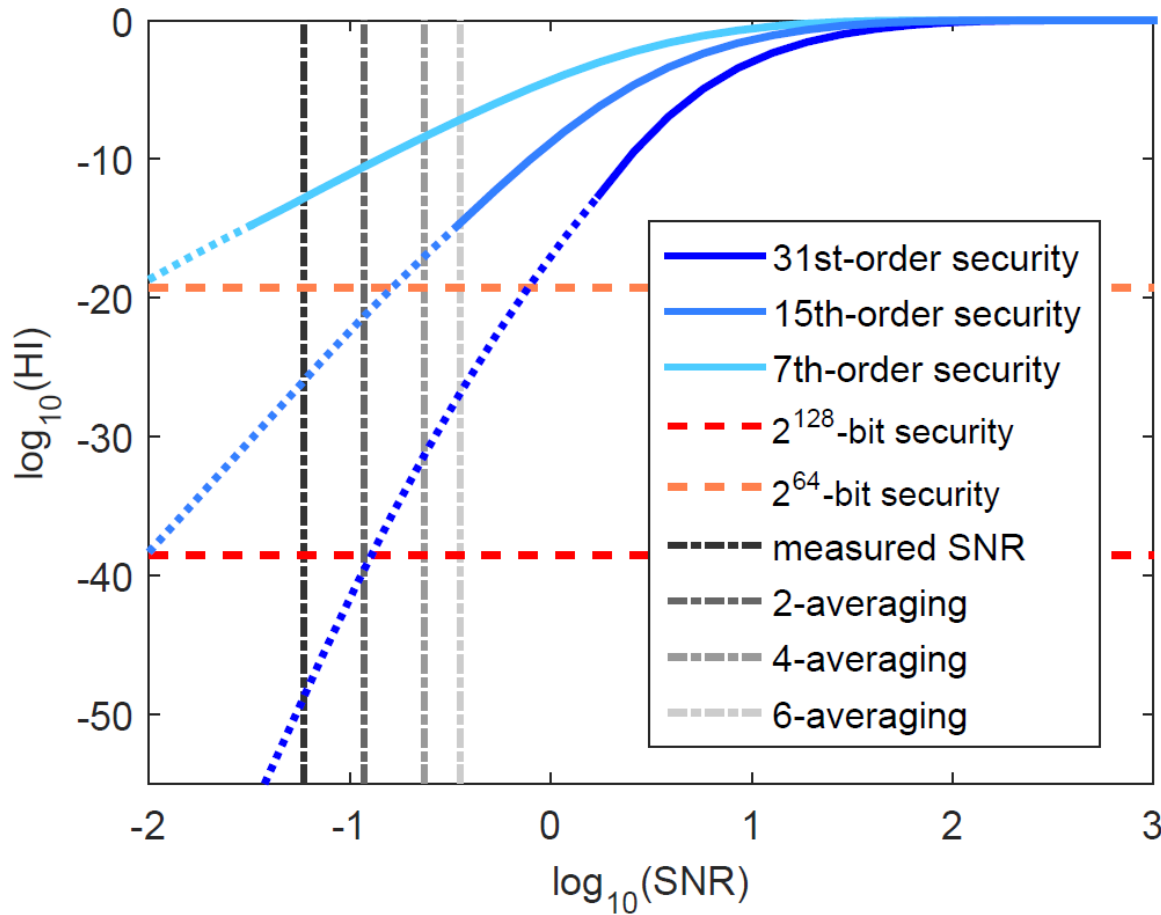




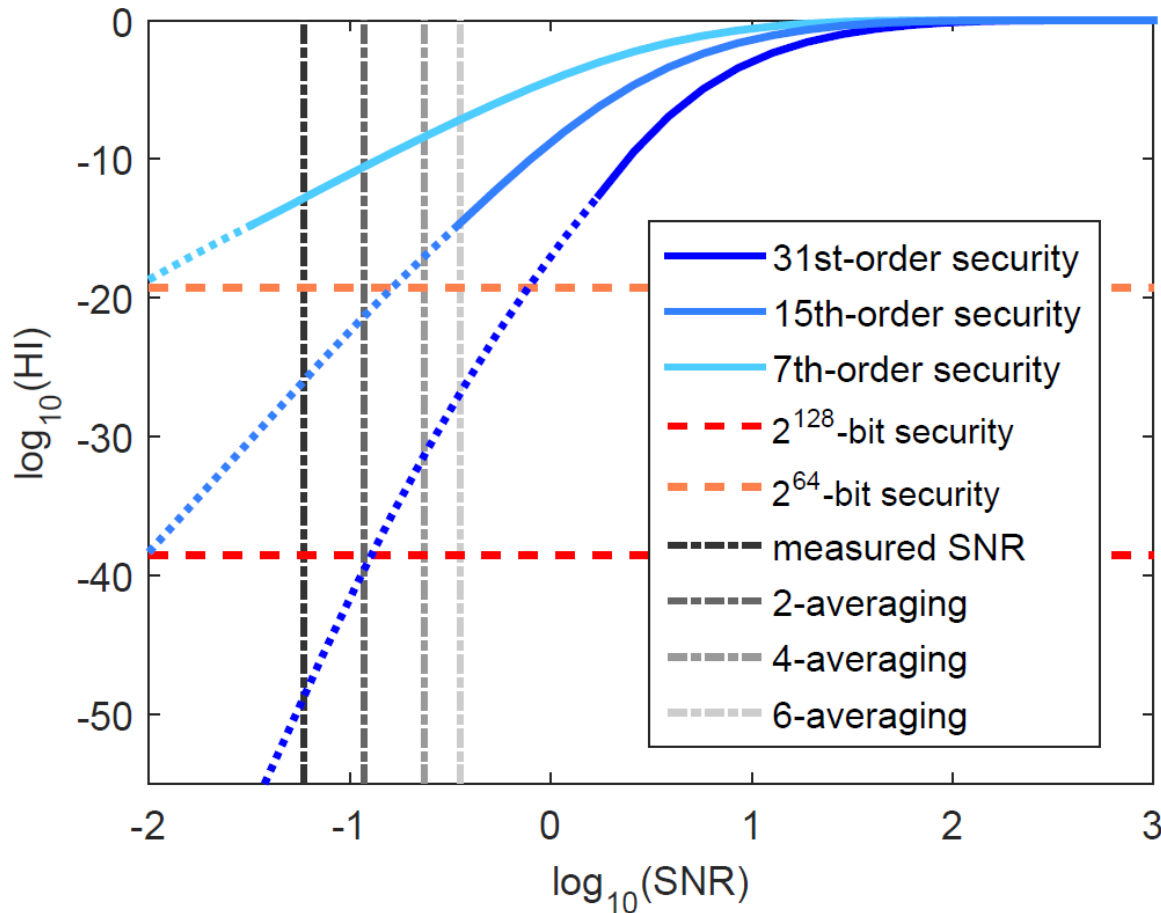
- SNR(=0,05) computed with linear regression
- **MI** of the encoding
- **31/15/7**-order security if flaw **f=1/2/4**



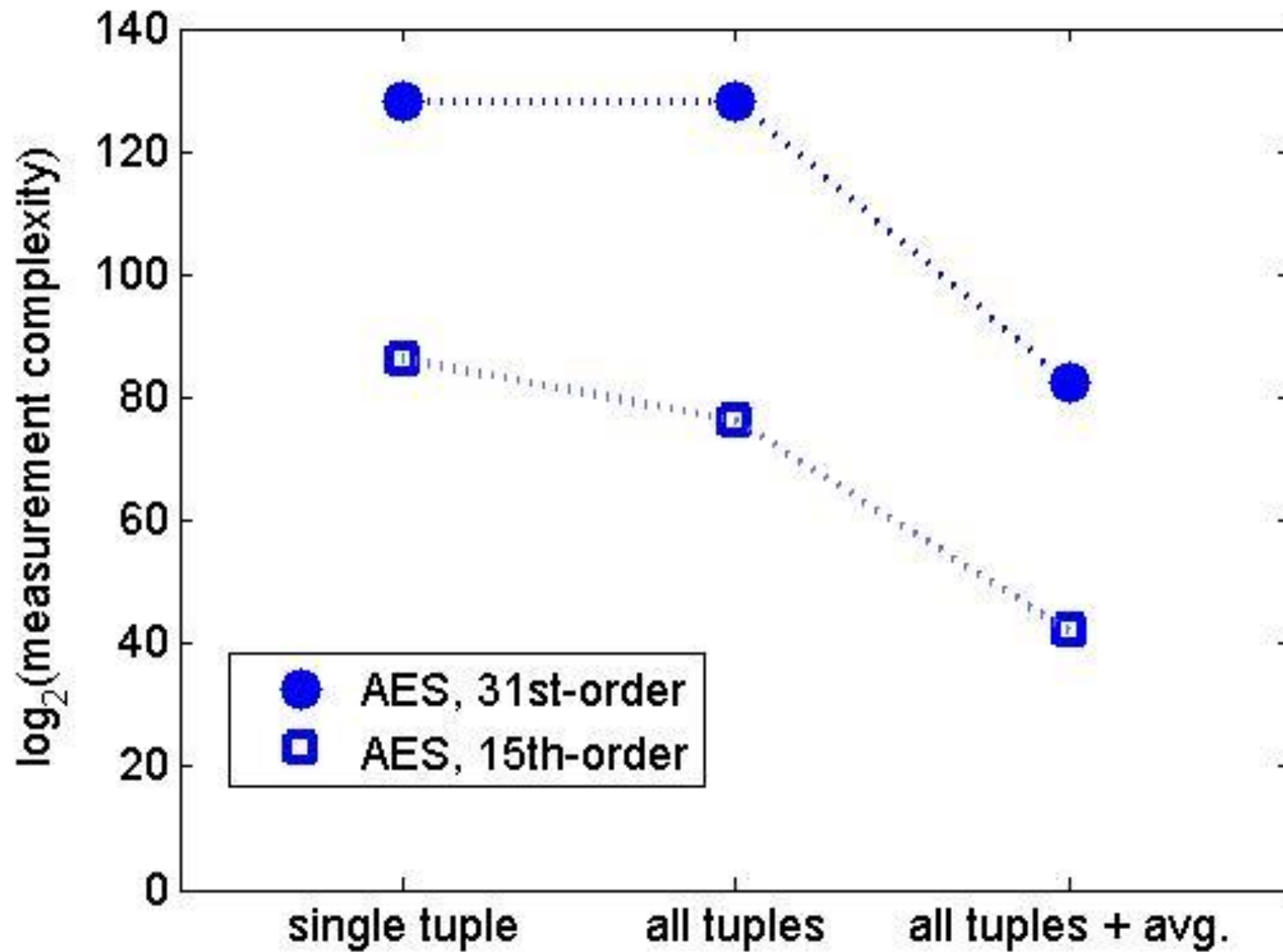
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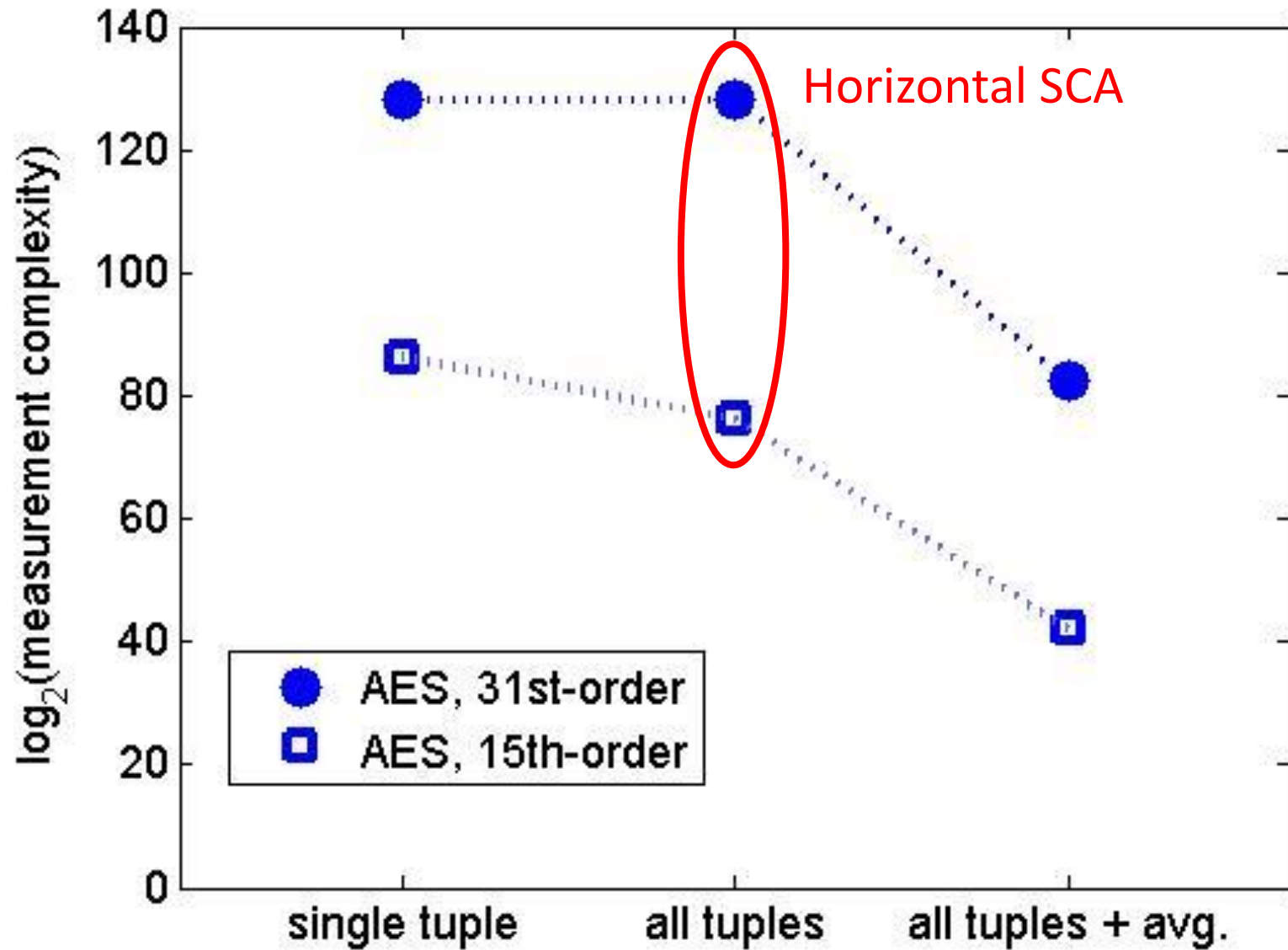


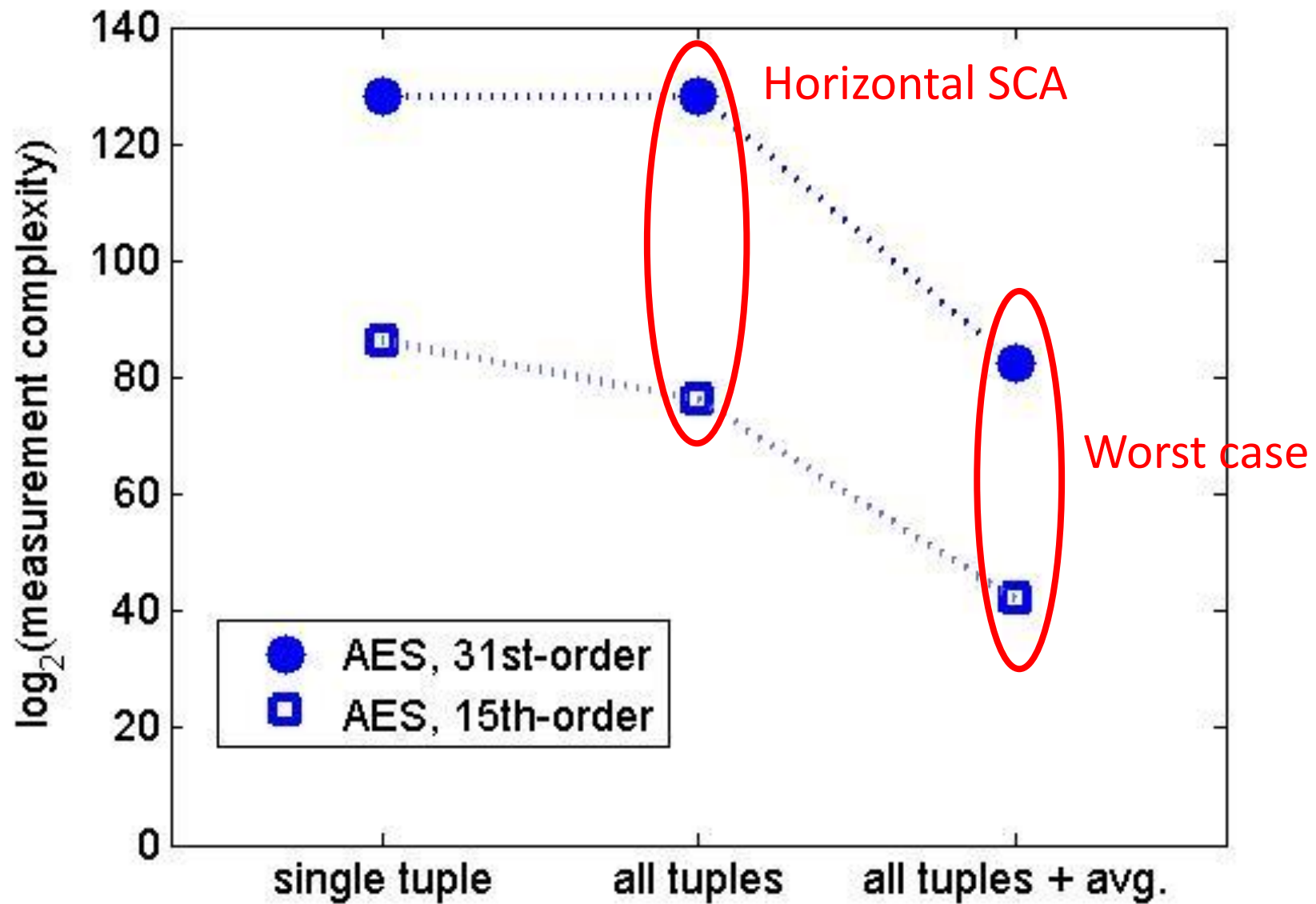
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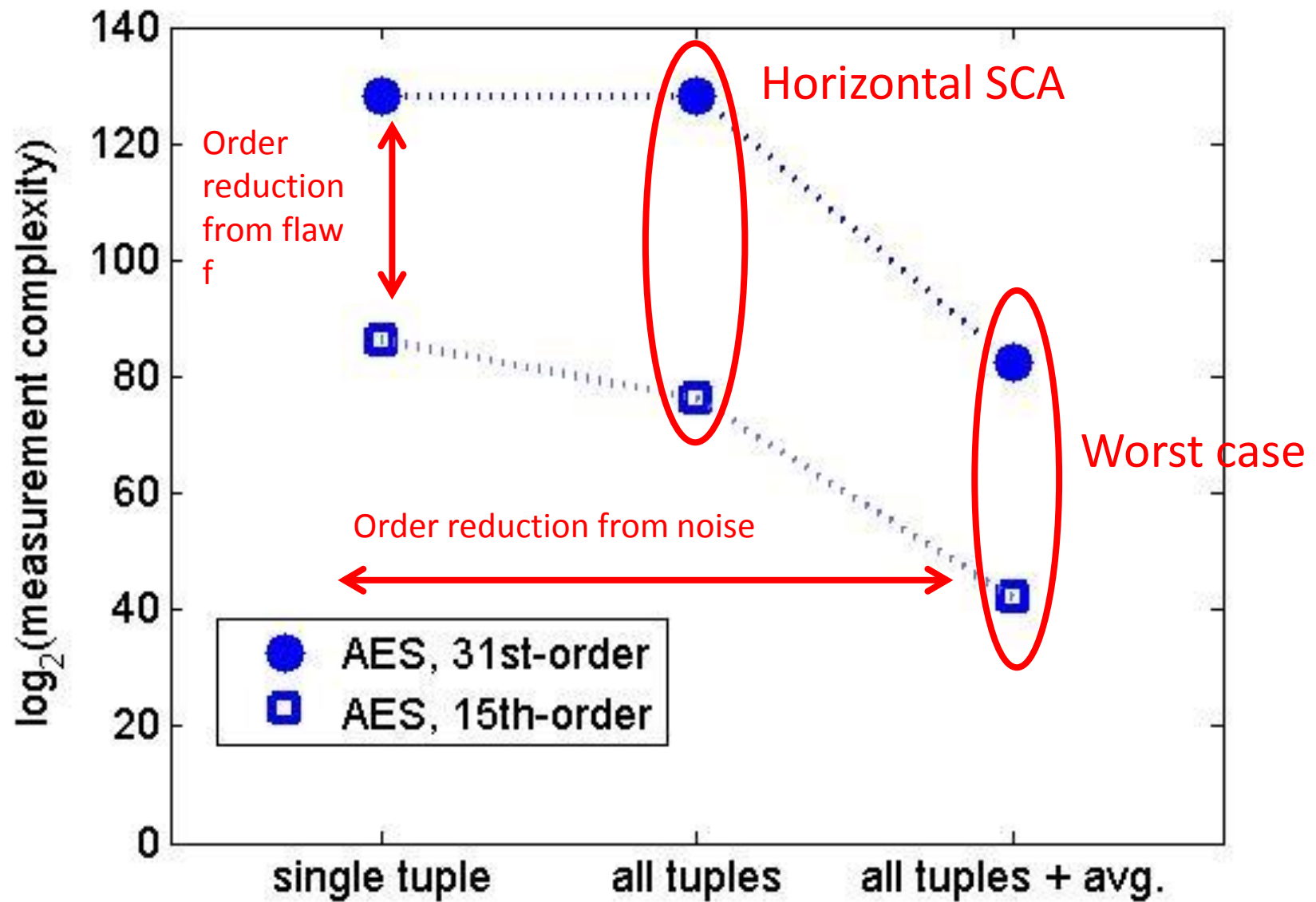


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- Averaging: multiple apparition of sensitive values









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 - Full code analysis to reduce refreshing
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**Thanks for your
attention**